

Gröbner
Bases Tutorial

David A. Cox

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Gröbner Bases Tutorial

Part I: Gröbner Bases and the Geometry of Elimination

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ISSAC 2007 Tutorial

Outline

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- k – **field** (often algebraically closed)
- $\mathbf{x}^\alpha = x_1^{\alpha_1} \cdots x_n^{\alpha_n}$ – **monomial** in x_1, \dots, x_n
- $c\mathbf{x}^\alpha$, $c \in k$ – **term** in x_1, \dots, x_n
- $k[\mathbf{x}] = k[x_1, \dots, x_n]$ – **polynomial ring** in n variables
- $\mathbb{A}^n = \mathbb{A}^n(k)$ – n -dimensional **affine space** over k
- $\mathbf{V}(I) = \mathbf{V}(f_1, \dots, f_s) \subseteq \mathbb{A}^n$ – **variety** of $I = \langle f_1, \dots, f_s \rangle$
- $\mathbf{I}(V) \subseteq k[\mathbf{x}]$ – **ideal** of the variety $V \subseteq \mathbb{A}^n$
- $\sqrt{I} = \{f \in k[\mathbf{x}] \mid \exists m f^m \in I\}$ – the **radical** of I

Recall that I is a **radical ideal** if $I = \sqrt{I}$.

Monomial Orders

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Definition

A **monomial order** is a total order $>$ on the set of monomials \mathbf{x}^α satisfying:

- $\mathbf{x}^\alpha > \mathbf{x}^\beta$ implies $\mathbf{x}^\alpha \mathbf{x}^\gamma > \mathbf{x}^\beta \mathbf{x}^\gamma$ for all \mathbf{x}^γ
- $\mathbf{x}^\alpha > 1$ for all $\mathbf{x}^\alpha \neq 1$

We often think of a monomial order as a total order on the set of exponent vectors $\alpha \in \mathbb{N}^n$.

Lemma

A monomial order is a well-ordering on the set of all monomials.

Examples of Monomial Orders

Examples

- **Lex order** with $x_1 > \dots > x_n$: $\mathbf{x}^\alpha >_{lex} \mathbf{x}^\beta$ iff

$$\alpha_1 > \beta_1, \text{ or } \alpha_1 = \beta_1 \text{ and } \alpha_2 > \beta_2, \text{ or } \dots$$

- **Weighted order** using $w \in \mathbb{R}_+^n$ and lex to break ties: $\mathbf{x}^\alpha >_{w,lex} \mathbf{x}^\beta$ iff

$$w \cdot \alpha > w \cdot \beta, \text{ or } w \cdot \alpha = w \cdot \beta \text{ and } \mathbf{x}^\alpha >_{lex} \mathbf{x}^\beta$$

- **Graded lex order** with $x_1 > \dots > x_n$: This is $>_{w,lex}$ for $w = (1, \dots, 1)$

Weighted orders will appear in Part II when we discuss the Gröbner walk.

Leading Terms

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Definition

Fix a monomial order $>$ and let $f \in k[\mathbf{x}]$ be nonzero. Write

$$f = c\mathbf{x}^\alpha + \text{terms with exponent vectors } \beta \neq \alpha$$

such that $c \neq 0$ and $\mathbf{x}^\alpha > \mathbf{x}^\beta$ wherever $\beta \neq \alpha$ and \mathbf{x}^β appears in a nonzero term of f . Then:

- $\text{LT}(f) = c\mathbf{x}^\alpha$ is the **leading term** of f
- $\text{LM}(f) = \mathbf{x}^\alpha$ is the **leading monomial** of f
- $\text{LC}(f) = c$ is the **leading coefficient** of f

The leading term $\text{LT}(f)$ is sometimes called the **initial term**, denoted $\text{in}(f)$.

The Division Algorithm

Division Algorithm

Given nonzero polynomials $f, f_1, \dots, f_s \in k[\mathbf{x}]$ and a monomial order $>$, there exist $r, q_1, \dots, q_s \in k[\mathbf{x}]$ with the following properties:

- $f = q_1 f_1 + \dots + q_s f_s + r.$
- No term of r is divisible by any of $\text{LT}(f_1), \dots, \text{LT}(f_s).$
- $\text{LT}(f) = \max_{>} \{ \text{LT}(q_i) \text{LT}(f_i) \mid q_i \neq 0 \}.$

Definition

Any representation

$$f = q_1 f_1 + \dots + q_s f_s$$

satisfying the third bullet is a **standard representation** of $f.$

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The Ideal of Leading Terms

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Definition

Given an ideal $I \subseteq k[\mathbf{x}]$ and a monomial order $>$, the **ideal of leading terms** is the monomial ideal

$$\langle \text{LT}(I) \rangle := \langle \text{LT}(f) \mid f \in I \rangle.$$

If $I = \langle f_1, \dots, f_s \rangle$, then

$$\langle \text{LT}(f_1), \dots, \text{LT}(f_s) \rangle \subseteq \langle \text{LT}(I) \rangle,$$

though equality **need not** occur.

This is where Gröbner bases enter the picture!

Gröbner Bases

Fix a monomial order $>$ on $k[\mathbf{x}]$.

Definition

Given an ideal $I \subseteq k[\mathbf{x}]$ a finite set $G \subseteq I \setminus \{0\}$ is a **Gröbner basis** for I under $>$ if

$$\langle \text{LT}(g) \mid g \in G \rangle = \langle \text{LT}(I) \rangle.$$

Definition

A Gröbner basis G is **reduced** if for every $g \in G$,

- $\text{LT}(g)$ divides no term of any element of $G \setminus \{g\}$.
- $\text{LC}(g) = 1$.

Theorem

Every ideal has a unique reduced Gröbner basis under $>$.

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Criteria to be a Gröbner Basis

Given $>$ and $g, h \in k[\mathbf{x}] \setminus \{0\}$, we get the **S-polynomial**

$$S(g, h) := \frac{\mathbf{x}^\gamma}{\text{LT}(g)} g - \frac{\mathbf{x}^\gamma}{\text{LT}(h)} h, \quad \mathbf{x}^\gamma = \text{lcm}(\text{LM}(g), \text{LM}(h)).$$

Three Criteria

- **(SR)** $G \subseteq I$ is a Gröbner basis of $I \iff$ every $f \in I$ has a standard representation using G .
- **(Buchberger)** G is a Gröbner basis of $\langle G \rangle \iff$ for every $g, h \in G$, $S(g, h)$ has a standard representation using G .
- **(LCM)** G is a Gröbner basis of $\langle G \rangle \iff$ for every $g, h \in G$, $S(g, h) = \sum_{\ell \in G} A_\ell \ell$, where $A_\ell \neq 0$ implies $\text{LT}(A_\ell \ell) < \text{lcm}(\text{LM}(g), \text{LM}(h))$ (a **lcm representation**).

The Consistency Theorem

Fix an ideal $I \subseteq k[\mathbf{x}]$, where k is algebraically closed.

Nullstellensatz

- **(Strong)** $\mathbf{I}(\mathbf{V}(I)) = \sqrt{I}$.
- **(Weak)** $\mathbf{V}(I) = \emptyset \iff 1 \in I \iff I = k[\mathbf{x}]$.

The Consistency Theorem

The following are equivalent:

- $I \neq k[\mathbf{x}]$.
- $1 \notin I$.
- $\mathbf{V}(I) \neq \emptyset$.
- I has a Gröbner basis consisting of nonconstant polynomials.
- I has a reduced Gröbner basis $\neq \{1\}$.

The Finiteness Theorem

Fix an ideal $I \subseteq k[\mathbf{x}]$, where k is algebraically closed. Also fix a monomial order $>$.

The Finiteness Theorem

The following are equivalent:

- $\mathbf{V}(I) \subseteq \mathbb{A}^n$ is finite.
- $k[\mathbf{x}]/I$ is a finite-dimensional vector space over k .
- I has a Gröbner basis G where $\forall i$, G has a element whose leading monomial is a power of x_i .
- Only finitely many monomials are not in $\langle \text{LT}(I) \rangle$.

When these conditions are satisfied:

- $\# \text{ solutions} \leq \dim_k k[\mathbf{x}]/I$.
- Equality holds $\iff I$ is radical.
- $\dim_k k[\mathbf{x}]/I = \# \text{ solutions counted with multiplicity}$.

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Given

$$k[\mathbf{x}, \mathbf{y}] = k[x_1, \dots, x_s, y_{s+1}, \dots, y_n],$$

we write monomials as $\mathbf{x}^\alpha \mathbf{y}^\beta$.

Definition

A monomial order $>$ on $k[\mathbf{x}, \mathbf{y}]$ **eliminates** \mathbf{x} whenever

$$\mathbf{x}^\alpha > \mathbf{x}^\beta \Rightarrow \mathbf{x}^\alpha \mathbf{y}^\gamma > \mathbf{x}^\beta \mathbf{y}^\delta$$

for all $\mathbf{y}^\gamma, \mathbf{y}^\delta$.

Example

Lex with $x_1 > \dots > x_n$ eliminates $\mathbf{x} = \{x_1, \dots, x_s\} \forall s$.

The Elimination Theorem

Fix an ideal $I \subseteq k[\mathbf{x}, \mathbf{y}]$.

Definition

$I \cap k[\mathbf{y}]$ is the **elimination ideal** of I that eliminates \mathbf{x} .

Theorem

Let G be a Gröbner basis of I for a monomial order $>$ that eliminates \mathbf{x} . Then $G \cap k[\mathbf{y}]$ is a Gröbner basis of the elimination ideal $I \cap k[\mathbf{y}]$ for the monomial order on $k[\mathbf{y}]$ induced by $>$.

Proof

$f \in I \cap k[\mathbf{y}]$ has standard representation $f = \sum_{g \in G} A_g g$. If $A_g \neq 0$, then $\text{LT}(g) \leq \text{LT}(A_g g) \leq \text{LT}(f) \in k[\mathbf{y}]$, so $g \in G \cap k[\mathbf{y}]$.
SR Criterion $\Rightarrow G \cap k[\mathbf{y}]$ is a Gröbner basis of $I \cap k[\mathbf{y}]$. \square

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Partial Solutions

Given $I \subseteq k[\mathbf{x}, \mathbf{y}] = k[x_1, \dots, x_s, y_{s+1}, \dots, y_n]$, the elimination ideal $I \cap k[\mathbf{y}]$ will be denoted

$$I_s := I \cap k[\mathbf{y}] \subseteq k[\mathbf{y}].$$

Definition

The **variety of partial solutions** is

$$\mathbf{V}(I_s) \subseteq \mathbb{A}^{n-s}.$$

Question

How do the partial solutions $\mathbf{V}(I_s) \subseteq \mathbb{A}^{n-s}$ relate to the original variety $V := \mathbf{V}(I) \subseteq \mathbb{A}^n$?

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Partial Solutions

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Given coordinates $x_1, \dots, x_s, y_{s+1}, \dots, y_n$, let

$$\pi_s : \mathbb{A}^n \longrightarrow \mathbb{A}^{n-s}$$

denote projection onto the last $n - s$ coordinates.

An ideal $I \subseteq k[\mathbf{x}, \mathbf{y}]$ gives:

- $V = \mathbf{V}(I) \subseteq \mathbb{A}^n$ and $\pi_s(V) \subseteq \mathbb{A}^{n-s}$.
- $I_s = I \cap k[\mathbf{y}] \subseteq k[\mathbf{y}]$ and $\mathbf{V}(I_s) \subseteq \mathbb{A}^{n-s}$.

Lemma

$$\pi_s(V) \subseteq \mathbf{V}(I_s).$$

Partial Solutions Don't Always Extend

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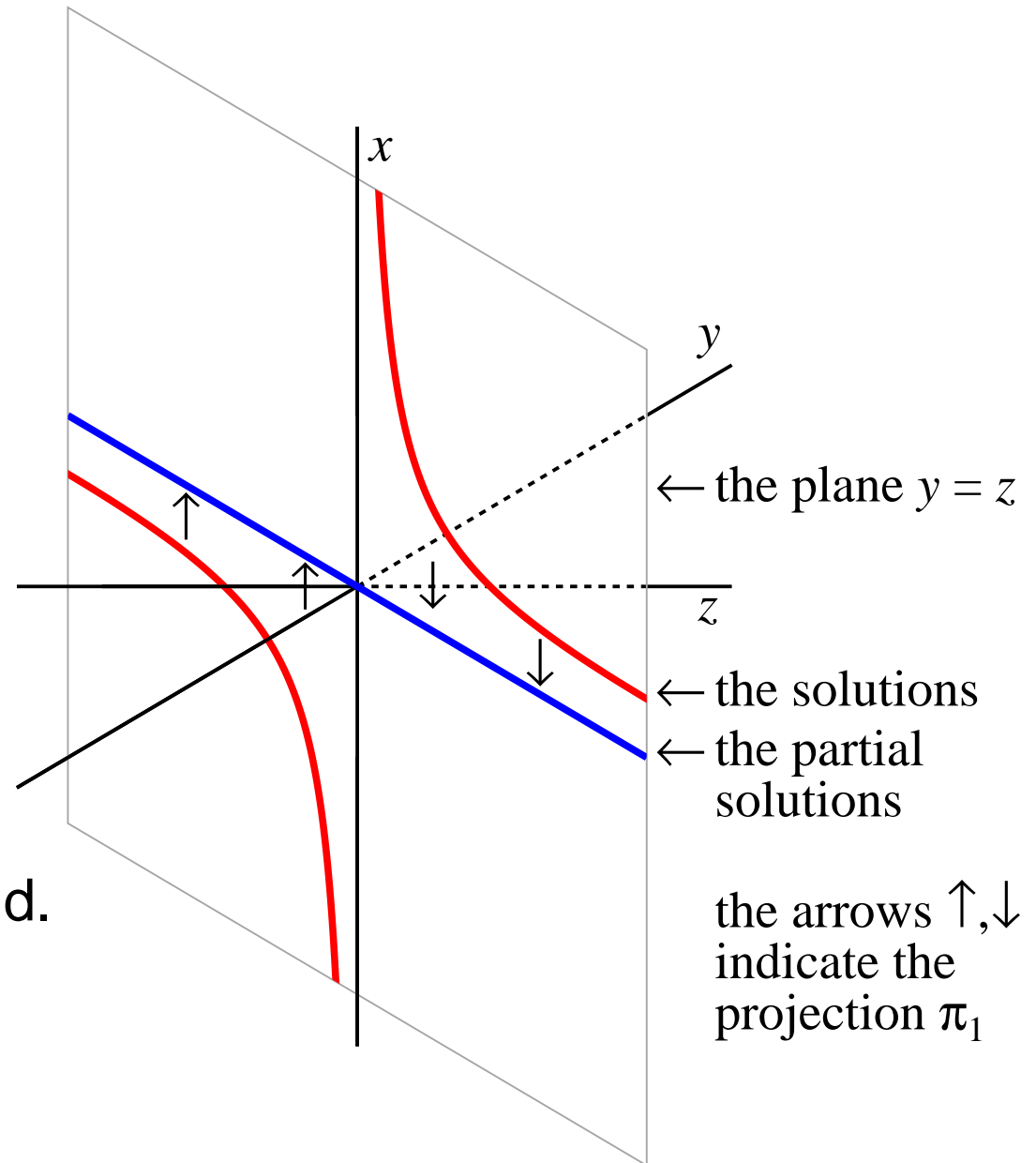
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In \mathbb{A}^3 , consider
 $V = V(xy - 1, y - z)$.
Using lex order
with $x > y > z$,
 $I = \langle xy - 1, y - z \rangle$
has Gröbner basis
 $xy - 1, y - z$. Thus
 $I_1 = \langle y - z \rangle$, so the
partial solutions are
the line $y = z$. The
partial solution
 $(0, 0)$ does **not** extend.



The Extension Theorem

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Let $I \subseteq k[x, y_2, \dots, y_n] = k[x, \mathbf{y}]$ with variety $V = \mathbf{V}(I) \subseteq \mathbb{A}^n$, and let $I_1 := I \cap k[\mathbf{y}]$ be the first elimination ideal. We assume that k is algebraically closed.

Theorem

Let $\mathbf{b} = (a_2, \dots, a_n) \in \mathbf{V}(I_1)$ be a partial solution. If the ideal I contains a polynomial f such that

$$f = c(\mathbf{y})x^N + \text{terms of degree} < N \text{ in } x$$

with $c(\mathbf{b}) \neq 0$, then there is $a \in k$ such that $(a, \mathbf{b}) = (a, a_2, \dots, a_n)$ is a solution, i.e.,

$$(a, a_2, \dots, a_n) \in V.$$

Zariski Closure

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Definition

Given a subset $S \subseteq \mathbb{A}^n$, the **Zariski closure** of S is the smallest variety $\overline{S} \subseteq \mathbb{A}^n$ containing S .

Lemma

The Zariski closure of $S \subseteq \mathbb{A}^n$ is $\overline{S} = \mathbf{V}(\mathbf{I}(S))$.

Example

Over \mathbb{C} , the set $\mathbb{Z}^n \subseteq \mathbb{C}^n$ has Zariski closure $\overline{\mathbb{Z}^n} = \mathbb{C}^n$.

The Closure Theorem

Let $V = \mathbf{V}(I) \subseteq \mathbb{A}^n$ and let k be algebraically closed.

Theorem

$$\mathbf{V}(I_s) = \overline{\pi_s(V)}.$$

*Thus $\mathbf{V}(I_s)$ is the smallest variety in \mathbb{A}^{n-s} containing $\pi_s(V)$.
Furthermore, there is an affine variety*

$$W \subseteq \mathbf{V}(I_s) \subseteq \mathbb{A}^{n-s}$$

with the following properties:

- $\overline{\mathbf{V}(I_s) \setminus W} = \mathbf{V}(I_s)$.
- $\mathbf{V}(I_s) \setminus W \subseteq \pi_s(V)$.

Thus “most” partial solutions in $\mathbf{V}(I_s)$ come from actual solutions, i.e, the projection of V fills up “most” of $\mathbf{V}(I_s)$.

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Traditional proofs of the Extension and Closure Theorems use resultants or more abstract methods from algebraic geometry.

Recently, Peter Schauenberg wrote

[“A Gröbner-based treatment of elimination theory for affine varieties”](#)

(Journal of Symbolic Computation, to appear). This paper uses Gröbner bases to give new proofs of the Extension and Closure Theorems.

Part I of the tutorial will conclude with these proofs. We begin with the Extension Theorem.

Prove the Extension Theorem

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References

Fix a partial solution $\mathbf{b} = (a_2, \dots, a_n) \in \mathbf{V}(I_1) \subset \mathbb{A}^{n-1}$.

Notation for the Proof

Let $f \in k[x, y_2, \dots, y_n] = k[x, \mathbf{y}]$.

- We write

$$f = \underbrace{c(\mathbf{y})}_{\text{LC}_x(f)} x^M + \text{terms of degree} < M \text{ in } x$$

- We set

$$\bar{f} := f(x, \mathbf{b}) \in k[x].$$

By hypothesis, there is $f \in I$ with $\overline{\text{LC}_x(f)} \neq 0$.

Lemma

Let G be a Gröbner basis of I using $>$ that eliminates x .

Lemma

There is $g \in G$ with $\overline{\text{LC}_x(g)} \neq 0$.

Proof

Let $f = \sum_{g \in G} A_g g$ be a standard representation of $f \in I$ with $\overline{\text{LC}_x(f)} \neq 0$. Thus

$$\text{LT}(f) = \max\{\text{LT}(A_g g) \mid A_g \neq 0\}.$$

Since $>$ eliminates x , it follows that

$$\deg_x(f) = \max\{\deg_x(A_g g) \mid A_g \neq 0\}$$

$$\text{LC}_x(f) = \sum_{\deg_x(A_g g) = \deg_x(f)} \text{LC}_x(A_g) \text{LC}_x(g)$$

□

Main Claim

By the lemma, we can pick $g \in G$ with $\overline{\text{LC}_x(g)} \neq 0$ and $M := \deg_x(g)$ minimal. Note that $M = \deg_x(\bar{g}) > 0$.

Main Claim

$$\{\bar{f} \mid f \in I\} = \langle \bar{g} \rangle \subseteq k[x].$$

Consequence: If $\bar{g}(a) = 0$ for some $a \in k$, then

$$f(a, \mathbf{b}) = \bar{f}(a) = 0$$

for all $f \in I$, so $(a, \mathbf{b}) = (a, a_2, \dots, a_n) \in V = \mathbf{V}(I)$.

This proves the Extension Theorem!

Strategy to prove Main Claim: Show $\bar{h} \in \langle \bar{g} \rangle$ for all $h \in G$.

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Claim

Consider $h \in G$ with $\deg_x(h) < M = \deg_x(g)$. M minimal implies $\overline{\text{LC}_x(h)} = 0$, so $\deg_x(\bar{h}) < \deg_x(h)$.

Claim

$\bar{h} = 0$.

Proof. Let $m := \deg_x(h) < M$ and set

$$S := \text{LC}_x(g)x^{M-m}h - \text{LC}_x(h)g \in I,$$

with standard representation $S = \sum_{\ell \in G} A_\ell \ell$. Observe:

- $\overline{\text{LC}_x(g)}x^{M-m}\bar{h} = \bar{S} = \sum_{\ell \in G} \bar{A}_\ell \bar{\ell}$
- $\max\{\deg_x(A_\ell) + \deg_x(\ell) \mid A_\ell \neq 0\} = \deg_x(S) < M$

Prove the Claim

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- $\overline{\text{LC}_x(g)} x^{M-m} \bar{h} = \bar{S} = \sum_{\ell \in G} \overline{A_\ell} \bar{\ell}$
- $\max\{\deg_x(A_\ell) + \deg_x(\ell)\} = \deg_x(S) < M$

First bullet implies: Since $\overline{\text{LC}_x(g)} \neq 0$ and $m = \deg_x(h)$,

$$M - \deg_x(h) + \deg_x(\bar{h}) \leq \max\{\deg_x(\overline{A_\ell}) + \deg_x(\bar{\ell})\}$$

so that

$$\deg_x(h) - \deg_x(\bar{h}) \geq \min\{M - (\deg_x(\overline{A_\ell}) + \deg_x(\bar{\ell}))\}.$$

Second bullet implies: All $\ell \in G$ in S have $\deg_x(\ell) < M$, so $\deg_x(\bar{\ell}) < \deg_x(\ell)$. Hence

$$\deg_x(\overline{A_\ell}) + \deg_x(\bar{\ell}) < \deg_x(A_\ell) + \deg_x(\ell) < M.$$

The two strict inequalities give $\deg_x(h) - \deg_x(\bar{h}) \geq 2$.

Finish the Claim

The inequality $\deg_x(h) - \deg_x(\bar{h}) \geq 2$ applies to **all** $h \in G$ with $\deg_x(h) < M$ and hence to **all** $\ell \in G$ in $S = \sum_{\ell} A_{\ell}\ell$.

Arguing as before gives

$$\deg_x(\overline{A_{\ell}}) + \deg_x(\bar{\ell}) < \deg_x(A_{\ell}) + \deg_x(\ell) < M,$$

\uparrow
drops by at least 2

and

$$\deg_x(h) - \deg_x(\bar{h}) \geq \underbrace{\min\{M - (\deg_x(\overline{A_{\ell}}) + \deg_x(\bar{\ell}))\}}_{\geq 3} \geq 3.$$

Continuing this way, we see that $\bar{h} = 0$ for all $h \in G$ with $\deg_x(h) < M$, as claimed. □

Prove the Main Claim

Proof. For $h \in G$, we show $\bar{h} \in \langle \bar{g} \rangle$ by induction on $\deg_x(h)$.

Base Case: $\deg_x(h) < M$ implies $\bar{h} = 0 \in \langle \bar{g} \rangle$ by Claim.

Inductive Step: Assume $\bar{h} \in \langle \bar{g} \rangle$ for all $h \in G$ with $\deg_x(h) < m \geq M$. Take $h \in G$ with $\deg_x(h) = m$. Then

$$S := \text{LC}_x(g)h - \text{LC}_x(h)x^{m-M}g \in I$$

has standard representation $S = \sum_{\ell \in G} A_\ell \ell$, so

$$\deg_x(S) < m \Rightarrow \deg_x(\ell) < m \quad \forall \ell \text{ in } S.$$

By inductive hypothesis, $\bar{\ell} \in \langle \bar{g} \rangle$. Hence

$$\overline{\text{LC}_x(g)h - \text{LC}_x(h)x^{m-M}g} = \bar{S} = \sum_{\ell \in G} \bar{A}_\ell \bar{\ell}.$$

Then $\overline{\text{LC}_x(g)} \neq 0$ implies $\bar{h} \in \langle \bar{g} \rangle$. □

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The Closure Theorem

We next give Schauenberg's proof of the Closure Theorem.

Fix a partial solution $\mathbf{b} = (a_{s+1}, \dots, a_n) \in \mathbf{V}(I_s) \subset \mathbb{A}^{n-s}$ and a Gröbner basis G of I for $>$ that eliminates $\mathbf{x} = (x_1, \dots, x_s)$.

Notation for the Proof

Let $f \in k[x_1, \dots, x_s, y_{s+1}, \dots, y_n] = k[\mathbf{x}, \mathbf{y}]$.

- We write

$$f = \underbrace{c(\mathbf{y})}_{\text{LC}_s(f)} \mathbf{x}^{\alpha(f)} + \text{terms} < \mathbf{x}^{\alpha(f)}.$$

- We set

$$\bar{f} := f(\mathbf{x}, \mathbf{b}) \in k[\mathbf{x}].$$

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A Special Case

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Special Case

If $\mathbf{b} \in \mathbf{V}(I_s)$ satisfies

$$\overline{\text{LC}_s(g)} \neq 0 \text{ for all } g \in G \setminus k[\mathbf{y}],$$

then $\mathbf{b} \in \pi_s(V)$.

Proof. If we can find $\mathbf{a} = (a_1, \dots, a_s)$ such that $\bar{g}(\mathbf{a}) = 0$ for all $g \in G$, then $g(\mathbf{a}, \mathbf{b}) = 0$ for all $g \in G$. This implies

$$(\mathbf{a}, \mathbf{b}) \in V,$$

and $\mathbf{b} \in \pi_s(V)$ follows.

Prove the Special Case

Let $\overline{G} = \{\bar{l} \mid l \in G \setminus k[\mathbf{y}]\}$. Take $g, h \in G \setminus k[\mathbf{y}]$ and set

$$S := \text{LC}_s(g) \frac{\mathbf{x}^\gamma}{\mathbf{x}^{\alpha(h)}} h - \text{LC}_s(h) \frac{\mathbf{x}^\gamma}{\mathbf{x}^{\alpha(g)}} g$$

where $\mathbf{x}^\gamma = \text{lcm}(\mathbf{x}^{\alpha(g)}, \mathbf{x}^{\alpha(h)})$. A standard representation $S = \sum_{l \in G} A_l l$ gives

$$\overline{\text{LC}_s(g)} \frac{\mathbf{x}^\gamma}{\mathbf{x}^{\alpha(h)}} \bar{h} - \overline{\text{LC}_s(h)} \frac{\mathbf{x}^\gamma}{\mathbf{x}^{\alpha(g)}} \bar{g} = \overline{S} = \sum_{\bar{l} \in \overline{G}} \overline{A}_l \bar{l}.$$

Then $\overline{\text{LC}_s(g)} \neq 0$ and $\overline{\text{LC}_s(h)} \neq 0$ imply:

- $\overline{S} = \sum_{\bar{l} \in \overline{G}} \overline{A}_l \bar{l}$ is a lcm representation.
- \overline{S} is the S-polynomial of \bar{g}, \bar{h} up to a constant.

Finish the Special Case

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- $\bar{S} = \sum_{\bar{\ell} \in \bar{G}} \bar{A}_{\bar{\ell}} \bar{\ell}$ is a lcm representation.
- \bar{S} is the S-polynomial of \bar{g}, \bar{h} up to a nonzero constant.

These tell us that for every

$$\bar{g}, \bar{h} \in \bar{G} = \{\bar{\ell} \mid \ell \in G \setminus k[\mathbf{y}]\},$$

the S-polynomial of \bar{g}, \bar{h} has a lcm representation with respect to \bar{G} . **LCM Criterion** $\Rightarrow \bar{G}$ is a Gröbner basis of $\langle \bar{G} \rangle$.

Since $\overline{\text{LT}_x(g)} \neq 0$ for $g \in G \setminus k[\mathbf{y}]$, \bar{g} is nonconstant for every $\bar{g} \in \bar{G}$. **Consistency Theorem** $\Rightarrow \mathbf{V}(\bar{G}) \neq \emptyset$.

Hence there exists \mathbf{a} such that $\bar{g}(\mathbf{a}) = 0$ for all $g \in G$. □

Saturation

Fix ideals $I, J \subseteq k[\mathbf{x}]$.

Definition

The **saturation** of I with respect to J is

$$I : J^\infty := \{f \in k[\mathbf{x}] \mid J^M f \subseteq I \text{ for } M \gg 0\}.$$

To see what this means geometrically, let

$$V := \mathbf{V}(I)$$

$$W := \mathbf{V}(J).$$

Then:

Lemma

$$\overline{V \setminus W} = \mathbf{V}(I : J^\infty).$$

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Proof of the Closure Theorem

Proof. The goal is to find a variety $W \subseteq \mathbf{V}(I_s)$ such that

$$\mathbf{V}(I_s) \setminus W \subseteq \pi_s(V) \quad \text{and} \quad \overline{\mathbf{V}(I_s) \setminus W} = \mathbf{V}(I_s).$$

Let G be a reduced Gröbner basis of I that eliminates \mathbf{x} . Set

$$J := I_s + \langle \prod_{g \in G \setminus k[\mathbf{y}]} \text{LC}_s(g) \rangle.$$

Then

$$\mathbf{V}(J) = \bigcup_{g \in G \setminus k[\mathbf{y}]} \mathbf{V}(I_s) \cap \mathbf{V}(\text{LC}_s(g)).$$

Notice that

$$\mathbf{b} \in \mathbf{V}(I_s) \setminus \mathbf{V}(J) \Rightarrow \overline{\text{LC}_s(g)} \neq 0 \quad \forall g \in G \setminus k[\mathbf{y}].$$

By **Special Case**, $\mathbf{b} \in \pi_s(V)$. Hence

$$\mathbf{V}(I_s) \setminus \mathbf{V}(J) \subseteq \pi_s(V).$$

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Two Cases

Case 1. $\overline{\mathbf{V}(I_s) \setminus \mathbf{V}(\text{LC}_s(g))} = \mathbf{V}(I_s)$ for **all** $g \in G \setminus k[\mathbf{y}]$.

Intersecting finitely many open dense sets is dense, so

$$\overline{\mathbf{V}(I_s) \setminus \mathbf{V}(J)} = \overline{\bigcap_{g \in G \setminus k[\mathbf{y}]} \mathbf{V}(I_s) \setminus \mathbf{V}(\text{LC}_s(g))} = \mathbf{V}(I_s).$$

Thus the theorem holds with $W = \mathbf{V}(J)$.

Case 2. $\overline{\mathbf{V}(I_s) \setminus \mathbf{V}(\text{LC}_s(g))} \subsetneq \mathbf{V}(I_s)$ for **some** $g \in G \setminus k[\mathbf{y}]$.

The strategy will be to enlarge I . First suppose that $\mathbf{V}(I_s) \subseteq \mathbf{V}(\text{LC}_s(g))$. Then:

- $\text{LC}_s(g)$ vanishes on $\mathbf{V}(I_s)$ and hence on V . Hence Nullstellensatz $\Rightarrow \text{LC}_s(g) \in \sqrt{I}$.
- G reduced $\Rightarrow \text{LC}_s(g) \notin I$ (since $\text{LT}(\text{LC}_s(g))$ divides $\text{LT}(g)$).

Finish the Proof

Together, these bullets imply that

$$I \subsetneq I + \langle \text{LC}_s(g) \rangle \subset \sqrt{I},$$

so in particular, $\mathbf{V}(I) = \mathbf{V}(I + \langle \text{LC}_s(g) \rangle)$. So it suffices to prove the theorem for $I + \langle \text{LC}_s(g) \rangle$ when $\mathbf{V}(I_s) \subseteq \mathbf{V}(\text{LC}_s(g))$.

On the other hand, when $\mathbf{V}(I_s) \not\subseteq \mathbf{V}(\text{LC}_s(g))$, we have a union of proper subsets

$$\begin{aligned} \mathbf{V}(I_s) &= \overline{\mathbf{V}(I_s) \setminus \mathbf{V}(\text{LC}_s(g))} \cup (\mathbf{V}(I_s) \cap \mathbf{V}(\langle \text{LC}_s(g) \rangle)) \\ &= \mathbf{V}(I_s : \text{LC}_s(g)^\infty) \cup \mathbf{V}(I_s + \langle \text{LC}_s(g) \rangle) \end{aligned}$$

Hence it suffices to prove the theorem for

$$I \subsetneq I + \langle I_s : \text{LC}_s(g)^\infty \rangle \quad \text{and} \quad I \subsetneq I + \langle \text{LC}_s(g) \rangle.$$

ACC \Rightarrow these enlargements occur only finitely often. □

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Example

Consider \mathbb{A}^4 with variables w, x, y, z . Let $\pi_1 : \mathbb{A}^4 \rightarrow \mathbb{A}^3$ be projection onto the last three coordinates. Set

$$f = x^5 + x^4 + 2 - yz$$

and $I = \langle (y - z)(fw - 1), (yw - 1)(fw - 1) \rangle$. This defines

$$V := \mathbf{V}(I) = \mathbf{V}(y - z, yw - 1) \cup \mathbf{V}((x^5 + x^4 + 2 - yz)w - 1) \subseteq \mathbb{A}^4.$$

Then:

- $\{g_1, g_2\} = \{(y - z)(fw - 1), (yw - 1)(fw - 1)\}$ is a Gröbner basis of I for lex with $w > x > y > z$.
- $I_1 = 0$.

Furthermore,

$$g_1 = ((y - z)(x^5 + x^4 + 2 - yz))w + \dots$$

$$g_2 = (y(x^5 + x^4 + 2 - yz))w^2 + \dots$$

Continue Example

As in the proof of the Closure Theorem, let

$$\begin{aligned} J &= I_1 + \langle \prod_{g \in G \setminus k[x,y,z]} \text{LC}_1(g) \rangle \\ &= \langle (y-z)(x^5 + x^4 + 2 - yz) \cdot y(x^5 + x^4 + 2 - yz) \rangle \\ &= \langle y(y-z)(x^5 + x^4 + 2 - yz)^2 \rangle. \end{aligned}$$

This satisfies Case 1 of the proof, so that

$$\mathbf{V}(I_1) \setminus \mathbf{V}(J) = \mathbb{A}^3 \setminus \mathbf{V}(J) \subseteq \pi_1(V).$$

However,

$$\begin{aligned} \pi_1(V) &= (\mathbb{A}^3 \setminus \mathbf{V}(x^5 + x^4 + 2 - yz)) \cup \\ &\quad (\mathbf{V}(y-z, x^5 + x^4 + 2 - yz) \setminus \mathbf{V}(x^5 + x^4 + 2, y, z)). \end{aligned}$$

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Definition

A set $S \subseteq \mathbb{A}^n$ is **constructible** if there are affine varieties $W_i \subseteq V_i \subseteq \mathbb{A}^n$, $i = 1, \dots, N$, such that

$$S = \bigcup_{i=1}^N (V_i \setminus W_i).$$

Theorem

Let k be algebraically closed and $\pi_s : \mathbb{A}^n \rightarrow \mathbb{A}^{n-s}$ be projection onto the last $n - s$ coordinates. If $V \subseteq \mathbb{A}^n$ is an affine variety, then $\pi_s(V) \subseteq \mathbb{A}^{n-s}$ is constructible.

The proof of the Closure Theorem can be adapted to give an **algorithm** for writing $\pi_s(V)$ as a constructible set.

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